

Expressiveness vs Complexity: the Case of a Logic for Reasoning about Time

HS

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1. Motivations

Formal logic is a breathtakingly versatile tool used in Artificial Intelligence (e.g., Knowledge Representation), mathematics, cognitive science, philosophy, among others, where by logic we mean a system which given premisses enables to infer their consequences. Two main properties of such systems are:

- *Expressive power* – what can be expressed in the language of a logic.

- *Computational complexity* – how much time/memory is needed to perform reasoning in a logic.

In general, **better expressiveness** has the price of **worse complexity**:

- Better Expressive Power

Better Computational Complexity —

3. HS Fragments

Full HS is *undecidable*, so we search for its better-behaved fragments. Let:



As a result, the *big questions* are:

- What reasoning may be performed (which logics are decidable)?
- What reasoning may be performed efficiently (which logics are in the complexity class P)?

My research refers these questions to the case of *Halpern-Shoham logic*.

2. Halpern-Shoham Logic (HS)

HS is a modal logic for *reasoning about temporal intervals*. Its modal operators enable us to access an interval which: current interval

- *begins* the current interval (B)
- proceeds *during* the current interval (D)
- ends the current interval (E)
- overlaps the current interval (O)
- *is adjacent to* the current interval (A)
- is later than the current interval (L)

or is in an inverse relation: B, D, E, O, A, L. Hence, the modal operators are:



horn $\varphi := \lambda \mid [\mathsf{U}](\lambda \wedge \ldots \wedge \lambda \to \lambda) \mid \varphi \wedge \varphi$ $\varphi := \lambda \mid [\mathsf{U}](\lambda \to \lambda) \mid [\mathsf{U}](\lambda \land \lambda \to \bot) \mid [\mathsf{U}](\lambda \lor \lambda) \mid \varphi \land \varphi$ krom $\varphi := \lambda \mid [\mathsf{U}](\lambda \to \lambda) \mid [\mathsf{U}](\lambda \land \lambda \to \bot) \mid \varphi \land \varphi$ core

- Where: -p is a propositional variable;
 - -i is a nominal (variable which is true in exactly one interval);
 - $-@_i\psi$ states that ψ holds in interval *i*;
 - $[U]\psi$ states that ψ holds in all intervals.

To obtain decidability we need to *disallow discrete and irreflexive time lines*.

► Particularly interesting is HS_{horn}^{\sqcup} which is expressive enough for applications to temporal databases and computationally efficient (in class P).

4. My Results

1. Augmenting HS_{horn}^{\sqcup} and HS_{core}^{\sqcup} with the ability of referring to single temporal intervals *does*

 $\langle \mathsf{R} \rangle \psi - \psi$ holds *in some* interval that is in relation R with the current one; $[\mathsf{R}]\psi - \psi$ holds in all intervals that are in relation R with the current one; where ψ is a formula, and $R \in \{B, \overline{B}, D, \overline{D}, E, \overline{E}, O, \overline{O}, A, \overline{A}, L, \overline{L}\}$.

Example: the formula $[L][L](conference \rightarrow \langle B \rangle opening \land \langle E \rangle closing)$ states that each conference is begun by an opening and ended by a closing.

not lead to undecidability.

2. The price for such referen*tiality* is (in most cases) *NP-completeness* (i.e., probably a loss of the efficiency of reasoning).



 $\mathsf{HS}_{horn}^{\Box,i,@}$

NP-complete

5. Overall Results – the Complexity Map of HS Fragments







6. Conclusions and Future Work

have constructed referential extensions of HS_{horn}^{\sqcup} and HS_{core}^{\sqcup} , and proved *their* NP*-completeness* (except $HS_{core}^{\sqcup,i}$), hence *decidability*.

Interesting *open questions* include the following:

 \triangleright What is the complexity of $HS_{core}^{\sqcup,i}$ (when discrete and irreflexive time lines) are disallowed)?

► Is HS_{core}^{\Diamond} decidable?

 \triangleright Which HS fragments allow referentiality without i and $@_i$?

7. References

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